

Two Results on Vector Gaussian Multi-Terminal Source Coding

Yinfei Xu

Email: xuyinfei@inc.cuhk.edu.hk

Joint work with Qiao Wang and Jun Chen

November 30, 2016



網絡編碼研究所
Institute of Network Coding

Outline

Two Results
on Vector
Gaussian
Multi-
Terminal
Source Coding

Yinfei Xu

Vector
Gaussian
Multiple
Description
Coding

Problem
Formulation
A Single-Letter
Lower Bound for
Sum Rate
Vector Gaussian
Case

Vector
Gaussian
One-Help-One
Problem

Problem
Formulation
KKT Conditions
The Extremal
Inequality
Rate Region

- 1 Vector Gaussian Multiple Description Coding
 - Problem Formulation
 - A Single-Letter Lower Bound for Sum Rate
 - Vector Gaussian Case

Outline

Two Results
on Vector
Gaussian
Multi-
Terminal
Source Coding

Yinfei Xu

Vector
Gaussian
Multiple
Description
Coding

Problem
Formulation
A Single-Letter
Lower Bound for
Sum Rate
Vector Gaussian
Case

Vector
Gaussian
One-Help-One
Problem

Problem
Formulation
KKT Conditions
The Extremal
Inequality
Rate Region

- 1 Vector Gaussian Multiple Description Coding
 - Problem Formulation
 - A Single-Letter Lower Bound for Sum Rate
 - Vector Gaussian Case
- 2 Vector Gaussian One-Help-One Problem
 - Problem Formulation
 - KKT Conditions
 - The Extremal Inequality
 - Rate Region

Multiple Description Coding

Two Results
on Vector
Gaussian
Multi-
Terminal
Source Coding

Yinfei Xu

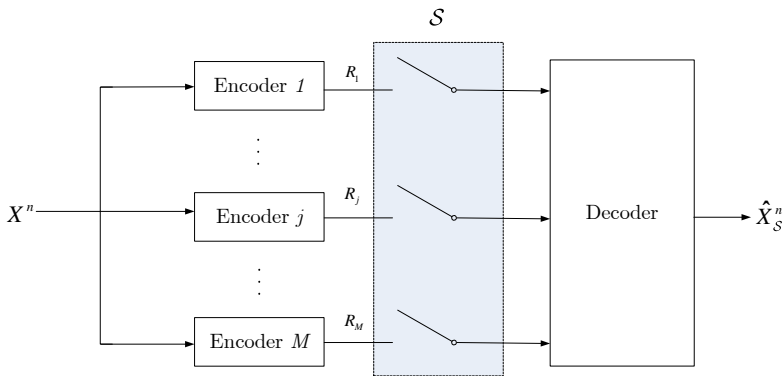
Vector
Gaussian
Multiple
Description
Coding

Problem
Formulation

A Single-Letter
Lower Bound for
Sum Rate
Vector Gaussian
Case

Vector
Gaussian
One-Help-One
Problem

Problem
Formulation
KKT Conditions
The Extremal
Inequality
Rate Region



Multiple Description Coding

Two Results on Vector Gaussian Multi-Terminal Source Coding

Yinfei Xu

Vector Gaussian Multiple Description Coding

Problem Formulation

A Single-Letter Lower Bound for Sum Rate
Vector Gaussian Case

Vector Gaussian One-Help-One Problem

Problem Formulation
KKT Conditions
The Extremal Inequality
Rate Region

- **Source:** i.i.d. process $\{X(t)\}_{t=1}^{\infty}$ with marginal distribution $p(x)$ over alphabet \mathcal{X} .
- **Encoding function:**

$$\varphi_i^{(n)} : \mathcal{X}^n \rightarrow \mathcal{C}_i^{(n)}, \quad i = 1, \dots, M,$$

$$\frac{1}{n} \log |\mathcal{C}_i^{(n)}| \leq R_i, \quad i = 1, \dots, M.$$

- **Decoding function:**

$$\psi_{\mathcal{S}}^{(n)} : \prod_{i \in \mathcal{S}} \mathcal{C}_i^{(n)} \rightarrow \hat{\mathcal{X}}^n, \quad \mathcal{S} \in 2_+^{\{1, \dots, M\}},$$

$$\frac{1}{n} \sum_{t=1}^n \mathbb{E}[d(X(t), \hat{X}_{\mathcal{S}}(t))] \leq d_{\mathcal{S}}, \quad \mathcal{S} \in 2_+^{\{1, \dots, M\}},$$

Tree-Structured Distortions

Two Results
on Vector
Gaussian
Multi-
Terminal
Source Coding

Yinfei Xu

Vector
Gaussian
Multiple
Description
Coding

Problem
Formulation

A Single-Letter
Lower Bound for
Sum Rate

Vector Gaussian
Case

Vector
Gaussian
One-Help-One
Problem

Problem
Formulation
KKT Conditions

The Extremal
Inequality
Rate Region

Definition

We say a set $\mathcal{T} \subseteq 2_+^{\{1, \dots, M\}}$ has a tree structure if, for any $\mathcal{S}_1, \mathcal{S}_2 \in \mathcal{T}$, one of the following statements is true:

- $\mathcal{S}_1 \subseteq \mathcal{S}_2$,
- $\mathcal{S}_2 \subseteq \mathcal{S}_1$,
- $\mathcal{S}_1 \cap \mathcal{S}_2 = \emptyset$.

Tree-Structured Distortions

Two Results
on Vector
Gaussian
Multi-
Terminal
Source Coding

Yinfei Xu

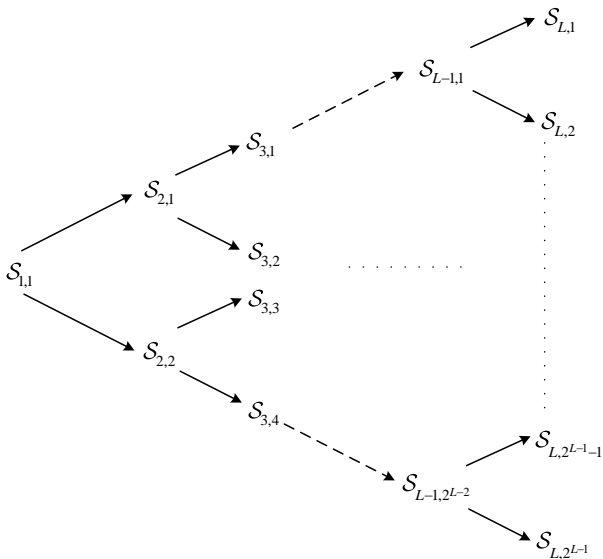
Vector
Gaussian
Multiple
Description
Coding

Problem
Formulation

A Single-Letter
Lower Bound for
Sum Rate
Vector Gaussian
Case

Vector
Gaussian
One-Help-One
Problem

Problem
Formulation
KKT Conditions
The Extremal
Inequality
Rate Region



Tree-Structured Distortions

Two Results
on Vector
Gaussian
Multi-
Terminal
Source Coding

Yinfei Xu

$$\mathcal{T} = \{S_{k,i} : k = 1, \dots, L; i = 1, \dots, 2^{k-1}\},$$
$$S_{k,i} = \left\{ j : \frac{2^L(i-1)}{2} < j \leq \frac{2^L i}{2^k} \right\}.$$

Vector
Gaussian
Multiple
Description
Coding

Problem
Formulation
A Single-Letter
Lower Bound for
Sum Rate
Vector Gaussian
Case

Vector
Gaussian
One-Help-One
Problem

Problem
Formulation
KKT Conditions
The Extremal
Inequality
Rate Region

- Two-description-three-distortion problem (L. Ozarow):
 $\mathcal{T} = \{\{1, 2\}, \{1\}, \{2\}\},$
- Individual and central distortion constraints (J. Chen):
 $\mathcal{T} = \{\{1, \dots, M\}, \{1\}, \dots, \{M\}\},$
- A **three-level-four-description** example:

$$\mathcal{T} = \{\{1, 2, 3, 4\}, \{1, 2\}, \{3, 4\}, \{1\}, \{2\}, \{3\}, \{4\}\}$$
$$= \{S_{1,1}, S_{2,1}, S_{2,2}, S_{3,1}, S_{3,2}, S_{3,3}, S_{3,4}\}$$

A Single-Letter Lower Bound for Sum Rate

- \mathcal{P} is the set of conditional distributions $p(\underline{z}|x)$ such that:

$$p(\underline{z}|x) = p(z_{1,1}|x)p(z_{2,1}, z_{2,2}|z_{1,1})$$

where $\underline{z} = (z_{1,1}, z_{2,1}, z_{2,2})$.

- $\mathcal{P}(\underline{d})$ is the set of conditional distributions $p(\hat{x}|x)$ such that the induced $p(x)p(\hat{x}|x)$ satisfies

$$\mathbb{E}[d(X, \hat{X}_{\mathcal{S}_{k,i}})] \leq d_{\mathcal{S}_{k,i}}, \quad k = 1, 2, 3; i = 1, \dots, 2^{k-1},$$

where $\underline{d} = (d_{\mathcal{S}_{k,i}})_{k=1,2,3;i=1,\dots,2^{k-1}}$ and
 $\hat{x} = (\hat{x}_{\mathcal{S}_{k,i}})_{k=1,2,3;i=1,\dots,2^{k-1}}$.

A Single-Letter Lower Bound for Sum Rate

Two Results
on Vector
Gaussian
Multi-
Terminal
Source Coding

Yinfei Xu

Vector
Gaussian
Multiple
Description
Coding

Problem
Formulation
A Single-Letter
Lower Bound for
Sum Rate
Vector Gaussian
Case

Vector
Gaussian
One-Help-One
Problem

Problem
Formulation
KKT Conditions
The Extremal
Inequality
Rate Region

Theorem

The sum rate $R(\underline{d})$ is defined as

$$R(\underline{d}) = \inf_{(R_1, R_2, R_3, R_4) \in \mathcal{R}(\underline{d})} R_1 + R_2 + R_3 + R_4,$$

which is lower bounded by

$$\begin{aligned} R(\underline{d}) \geq & \inf_{p(\hat{x}|x) \in \mathcal{P}(\underline{d})} \sup_{p(\underline{z}|x) \in \mathcal{P}} I(X; \hat{X}_{\mathcal{S}_{1,1}} | Z_{1,1}) \\ & + I(Z_{1,1}; \hat{X}_{\mathcal{S}_{2,1}} | Z_{2,1}) + I(Z_{1,1}; \hat{X}_{\mathcal{S}_{2,2}} | Z_{2,2}) \\ & + I(Z_{2,1}; \hat{X}_{\mathcal{S}_{3,1}}) + I(Z_{2,1}; \hat{X}_{\mathcal{S}_{3,2}}) \\ & + I(Z_{2,2}; \hat{X}_{\mathcal{S}_{3,3}}) + I(Z_{2,2}; \hat{X}_{\mathcal{S}_{3,4}}). \end{aligned}$$

where $p(\underline{z}, x, \hat{x}) = p(\underline{z}|x)p(x)p(\hat{x}|x)$.

Proof

Two Results on Vector Gaussian Multi- Terminal Source Coding

Yinfei Xu

Vector
Gaussian
Multiple
Description
Coding
Problem
Formulation
A Single-Letter
Lower Bound for
Sum Rate
Vector Gaussian
Case

Vector
Gaussian
One-Help-One
Problem
Problem
Formulation
KKT Conditions
The Extremal
Inequality
Rate Region

$$\begin{aligned} & \log |\mathcal{C}_1^{(n)}| + \log |\mathcal{C}_2^{(n)}| + \log |\mathcal{C}_3^{(n)}| + \log |\mathcal{C}_4^{(n)}| \\ & \geq H(\varphi_1^{(n)}) + H(\varphi_2^{(n)}) + H(\varphi_3^{(n)}) + H(\varphi_4^{(n)}) \\ & = H(\varphi_1^{(n)}, \varphi_2^{(n)}) + H(\varphi_3^{(n)}, \varphi_4^{(n)}) + I(\varphi_1^{(n)}; \varphi_2^{(n)}) + I(\varphi_3^{(n)}; \varphi_4^{(n)}) \\ & = H(\varphi_1^{(n)}, \varphi_2^{(n)}, \varphi_3^{(n)}, \varphi_4^{(n)}) + I(\varphi_1^{(n)}; \varphi_2^{(n)}) + I(\varphi_3^{(n)}; \varphi_4^{(n)}) \\ & \quad + I(\varphi_1^{(n)}, \varphi_2^{(n)}; \varphi_3^{(n)}, \varphi_4^{(n)}) \end{aligned}$$

- Introducing auxiliary remote source $Z_{2,1}^n$

$$\begin{aligned}
 & I(\varphi_1^{(n)}; \varphi_2^{(n)}) \\
 &= I(Z_{2,1}^n; \varphi_1^n) + I(Z_{2,1}^n; \varphi_2^{(n)}) - I(Z_{2,1}^n; \varphi_1^{(n)}, \varphi_2^{(n)}) \\
 &\quad + I(\varphi_1^n; \varphi_2^{(n)} | Z_{2,1}^n) \\
 &\geq I(Z_{2,1}^n; \varphi_1^n) + I(Z_{2,1}^n; \varphi_2^{(n)}) - I(Z_{2,1}^n; \varphi_1^{(n)}, \varphi_2^{(n)}).
 \end{aligned}$$

- Similarly introducing remote sources

$$\begin{aligned}
 & Z_{2,2}^n \text{ to } I(\varphi_3^{(n)}; \varphi_4^{(n)}), \\
 & Z_{1,1}^n \text{ to } I(\varphi_1^{(n)}, \varphi_2^{(n)}; \varphi_3^{(n)}, \varphi_4^{(n)}).
 \end{aligned}$$

Proof

Two Results
on Vector
Gaussian
Multi-
Terminal
Source Coding

Yinfei Xu

Vector
Gaussian
Multiple
Description
Coding

Problem
Formulation
A Single-Letter
Lower Bound for
Sum Rate
Vector Gaussian
Case

Vector
Gaussian
One-Help-One
Problem

Problem
Formulation
KKT Conditions
The Extremal
Inequality
Rate Region

$$\begin{aligned} & \log |\mathcal{C}_1^{(n)}| + \log |\mathcal{C}_2^{(n)}| + \log |\mathcal{C}_3^{(n)}| + \log |\mathcal{C}_4^{(n)}| \\ & \geq I(X^n; \varphi_1^{(n)}, \varphi_2^{(n)}, \varphi_3^{(n)}, \varphi_4^{(n)}) - I(\mathbf{Z}_{1,1}^n; \varphi_1^{(n)}, \varphi_2^{(n)}, \varphi_3^{(n)}, \varphi_4^{(n)}) \\ & \quad + I(\mathbf{Z}_{1,1}^n; \varphi_1^{(n)}, \varphi_2^{(n)}) - I(\mathbf{Z}_{2,1}^n; \varphi_1^{(n)}, \varphi_2^{(n)}) \\ & \quad + I(\mathbf{Z}_{1,1}^n; \varphi_3^{(n)}, \varphi_4^{(n)}) - I(\mathbf{Z}_{2,2}^n; \varphi_3^{(n)}, \varphi_4^{(n)}) \\ & \quad + I(\mathbf{Z}_{2,1}^n; \varphi_1^{(n)}) \\ & \quad + I(\mathbf{Z}_{2,1}^n; \varphi_2^{(n)}) \\ & \quad + I(\mathbf{Z}_{2,2}^n; \varphi_3^{(n)}) \\ & \quad + I(\mathbf{Z}_{2,2}^n; \varphi_4^{(n)}) \end{aligned}$$

Proof

Two Results
on Vector
Gaussian
Multi-
Terminal
Source Coding

Yinfei Xu

Vector
Gaussian
Multiple
Description
Coding

Problem
Formulation
A Single-Letter
Lower Bound for
Sum Rate
Vector Gaussian
Case

Vector
Gaussian
One-Help-One
Problem

Problem
Formulation
KKT Conditions
The Extremal
Inequality
Rate Region

- Single letterization of each term by time sharing random variable: T

e.g.,

$$\begin{aligned} & I(\mathbf{Z}_{1,1}^n; \varphi_1^{(n)}, \varphi_2^{(n)}) - I(\mathbf{Z}_{2,1}^n; \varphi_1^{(n)}, \varphi_2^{(n)}) \\ &= I(\mathbf{Z}_{1,1}^n; \varphi_1^{(n)}, \varphi_2^{(n)} | \mathbf{Z}_{2,1}^n) \quad (X^n \rightarrow \mathbf{Z}_{1,1}^n \rightarrow \mathbf{Z}_{2,1}^n) \\ &\geq I(\mathbf{Z}_{1,1}^n; \hat{X}_{\mathcal{S}_{2,1}}^n | \mathbf{Z}_{2,1}^n) \\ &\geq nI(\mathbf{Z}_{1,1}(T); \hat{X}_{\mathcal{S}_{1,1}}(T) | \mathbf{Z}_{2,1}(T)). \end{aligned}$$

Vector Gaussian Case: Converse Theorem

Two Results
on Vector
Gaussian
Multi-
Terminal
Source Coding

Yinfei Xu

Vector
Gaussian
Multiple
Description
Coding

Problem
Formulation
A Single-Letter
Lower Bound for
Sum Rate
Vector Gaussian
Case

Vector
Gaussian
One-Help-One
Problem

Problem
Formulation
KKT Conditions
The Extremal
Inequality
Rate Region

- Construction of auxiliary remote sources:

$$Z_{1,1} = X + N_{1,1}, \quad Z_{2,1} = X + N_{2,1}, \quad Z_{2,2} = X + N_{2,2},$$

where

$$\mathbf{0} \preceq \Sigma_{N_{1,1}} \preceq \Sigma_{N_{2,1}}, \Sigma_{N_{2,2}}.$$

- Introducing notations

$$\Theta_{1,1} = \left(\Sigma_X^{-1} + \Sigma_{N_{1,1}}^{-1} \right)^{-1},$$

$$\Theta_{2,1} = \left(\Sigma_X^{-1} + \Sigma_{N_{2,1}}^{-1} \right)^{-1}, \quad \Theta_{2,2} = \left(\Sigma_X^{-1} + \Sigma_{N_{2,2}}^{-1} \right)^{-1}.$$

$$\Rightarrow \Sigma_X \succeq \Theta_{2,1}, \Theta_{2,2} \succeq \Theta_{1,1} \succeq \mathbf{0}.$$

Vector Gaussian Case: Converse Theorem

Two Results
on Vector
Gaussian
Multi-
Terminal
Source Coding

Yinfei Xu

Vector
Gaussian
Multiple
Description
Coding

Problem
Formulation
A Single-Letter
Lower Bound for
Sum Rate
Vector Gaussian
Case

Vector
Gaussian
One-Help-One
Problem

Problem
Formulation
KKT Conditions
The Extremal
Inequality
Rate Region

- Applying worst noise lemma to each term: e.g.

$$\begin{aligned} & I(X; \hat{X}_{\mathcal{S}_{1,1}} | Z_{1,1}) \\ &= I(X; \hat{X}_{\mathcal{S}_{1,1}}) - I(Z_{1,1}; \hat{X}_{\mathcal{S}_{1,1}}) \\ &= \frac{1}{2} \log \left(\frac{\Sigma_X}{\Sigma_X + \Sigma_{N_{1,1}}} \right) + I(N_{1,1}; X + N_{1,1} | \hat{X}_{\mathcal{S}_{1,1}}) \\ &\geq \frac{1}{2} \log \frac{|\Sigma_X| |\mathbf{D}_{\mathcal{S}_{1,1}} + \Sigma_{N_{1,1}}|}{|\mathbf{D}_{\mathcal{S}_{1,1}}| |\Sigma_X + \Sigma_{N_{1,1}}|} \\ &= \frac{1}{2} \log \frac{|\Theta_{1,1}^{-1} + \mathbf{D}_{\mathcal{S}_{1,1}}^{-1} - \Sigma_X^{-1}|}{|\Theta_{1,1}^{-1}|} \end{aligned}$$

Vector Gaussian Case: Converse Theorem

Two Results
on Vector
Gaussian
Multi-
Terminal
Source Coding

Yinfei Xu

Vector
Gaussian
Multiple
Description
Coding

Problem
Formulation
A Single-Letter
Lower Bound for
Sum Rate
Vector Gaussian
Case

Vector
Gaussian
One-Help-One
Problem

Problem
Formulation
KKT Conditions
The Extremal
Inequality
Rate Region

$$\begin{aligned} R(\underline{d}) &\geq \max_{\Theta_{1,1}, \Theta_{2,1}, \Theta_{2,2}} \frac{1}{2} \log \frac{|\Theta_{1,1}^{-1} + \mathbf{D}_{S_{1,1}}^{-1} - \Sigma_X^{-1}|}{|\Theta_{1,1}^{-1}|} \\ &\quad + \frac{1}{2} \log \frac{|\Theta_{1,1}^{-1}| |\Theta_{2,1}^{-1} + \mathbf{D}_{S_{2,1}}^{-1} - \Sigma_X^{-1}|}{|\Theta_{2,1}^{-1}| |\Theta_{2,1}^{-1} + \mathbf{D}_{S_{2,1}}^{-1} - \Sigma_X^{-1}|} \\ &\quad + \frac{1}{2} \log \frac{|\Theta_{1,1}^{-1}| |\Theta_{2,2}^{-1} + \mathbf{D}_{S_{2,2}}^{-1} - \Sigma_X^{-1}|}{|\Theta_{2,2}^{-1}| |\Theta_{2,2}^{-1} + \mathbf{D}_{S_{2,2}}^{-1} - \Sigma_X^{-1}|} \\ &\quad + \sum_{i=1}^2 \frac{1}{2} \log \frac{|\Theta_{2,1}^{-1}| |\mathbf{D}_{S_{3,i}}^{-1}|}{|\Sigma_X^{-1}| |\Theta_{2,1}^{-1} + \mathbf{D}_{S_{3,i}}^{-1} - \Sigma_X^{-1}|} \\ &\quad + \sum_{i=3}^4 \frac{1}{2} \log \frac{|\Theta_{2,2}^{-1}| |\mathbf{D}_{S_{3,i}}^{-1}|}{|\Sigma_X^{-1}| |\Theta_{2,2}^{-1} + \mathbf{D}_{S_{3,i}}^{-1} - \Sigma_X^{-1}|} \\ &\text{subject to } \Sigma_X \succeq \Theta_{2,1}, \Theta_{2,2} \succeq \Theta_{1,1} \succeq \mathbf{0} \end{aligned}$$

Vector Gaussian Case: Direct Theorem

Two Results
on Vector
Gaussian
Multi-
Terminal
Source Coding

Yinfei Xu

Vector
Gaussian
Multiple
Description
Coding

Problem
Formulation
A Single-Letter
Lower Bound for
Sum Rate
Vector Gaussian
Case

Vector
Gaussian
One-Help-One
Problem

Problem
Formulation
KKT Conditions
The Extremal
Inequality
Rate Region

$$\Sigma_{\mathcal{S}_{k,i}} \triangleq (\mathbf{D}_{\mathcal{S}_{k,i}}^{-1} - \Sigma_X^{-1})^{-1} \succ \mathbf{0}, \quad k = 1, 2, 3; i = 1, \dots, 2^{k-1}.$$

$$\begin{aligned} & \max_{\underline{\Theta}} \frac{1}{2} \log \frac{|\Sigma_X + \Sigma_{\mathcal{S}_{1,1}}|}{|\Sigma_{\mathcal{S}_{1,1}}|} \\ & + \frac{1}{2} \log \frac{|\Theta_{1,1} + \Sigma_{\mathcal{S}_{1,1}}| |\Sigma_X + \Sigma_{\mathcal{S}_{2,1}}| |\Sigma_X + \Sigma_{\mathcal{S}_{2,2}}|}{|\Sigma_X + \Sigma_{\mathcal{S}_{1,1}}| |\Theta_{1,1} + \Sigma_{\mathcal{S}_{2,1}}| |\Theta_{1,1} + \Sigma_{\mathcal{S}_{2,2}}|} \\ & + \frac{1}{2} \log \frac{|\Theta_{2,1} + \Sigma_{\mathcal{S}_{2,1}}| |\Sigma_X + \Sigma_{\mathcal{S}_{3,1}}| |\Sigma_X + \Sigma_{\mathcal{S}_{3,2}}|}{|\Sigma_X + \Sigma_{\mathcal{S}_{2,1}}| |\Theta_{2,1} + \Sigma_{\mathcal{S}_{3,1}}| |\Theta_{2,1} + \Sigma_{\mathcal{S}_{3,2}}|} \\ & + \frac{1}{2} \log \frac{|\Theta_{2,2} + \Sigma_{\mathcal{S}_{2,2}}| |\Sigma_X + \Sigma_{\mathcal{S}_{3,3}}| |\Sigma_X + \Sigma_{\mathcal{S}_{3,4}}|}{|\Sigma_X + \Sigma_{\mathcal{S}_{2,2}}| |\Theta_{2,2} + \Sigma_{\mathcal{S}_{3,3}}| |\Theta_{2,2} + \Sigma_{\mathcal{S}_{3,4}}|} \end{aligned}$$

$$\text{subject to } \Sigma_X \succeq \Theta_{2,1}, \Theta_{2,2} \succeq \Theta_{1,1} \succeq \mathbf{0}$$

Vector Gaussian Case: Direct Theorem

Two Results
on Vector
Gaussian
Multi-
Terminal
Source Coding

Yinfei Xu

Vector
Gaussian
Multiple
Description
Coding

Problem
Formulation
A Single-Letter
Lower Bound for
Sum Rate
Vector Gaussian
Case

Vector
Gaussian
One-Help-One
Problem

Problem
Formulation
KKT Conditions
The Extremal
Inequality
Rate Region

KKT conditions

Let $(\Theta_{1,1}^*, \Theta_{2,1}^*, \Theta_{2,2}^*)$ be an optimizer, then there exists $(\mathbf{M}_{1,1}, \mathbf{M}_{2,1}, \mathbf{M}_{2,2}, \mathbf{M}_{3,1}, \mathbf{M}_{3,2})$ such that:

$$\begin{aligned} (\Theta_{1,1}^* + \Sigma_{S_{2,1}})^{-1} + \mathbf{M}_{2,1} + (\Theta_{1,1}^* + \Sigma_{S_{2,2}})^{-1} + \mathbf{M}_{2,2} \\ = (\Theta_{1,1}^* + \Sigma_{S_{1,1}})^{-1} + \mathbf{M}_{1,1} \end{aligned} \quad (1.1)$$

$$\begin{aligned} (\Theta_{2,1}^* + \Sigma_{S_{3,1}})^{-1} + \mathbf{M}_{3,1} + (\Theta_{2,1}^* + \Sigma_{S_{3,2}})^{-1} \\ = (\Theta_{2,1}^* + \Sigma_{S_{2,1}})^{-1} + \mathbf{M}_{2,1} \end{aligned} \quad (1.2)$$

$$\begin{aligned} (\Theta_{2,2}^* + \Sigma_{S_{3,3}})^{-1} + \mathbf{M}_{3,2} + (\Theta_{2,2}^* + \Sigma_{S_{3,4}})^{-1} \\ = (\Theta_{2,2}^* + \Sigma_{S_{2,2}})^{-1} + \mathbf{M}_{2,2}, \end{aligned} \quad (1.3)$$

Vector Gaussian Case: Direct Theorem

Two Results
on Vector
Gaussian
Multi-
Terminal
Source Coding

Yinfei Xu

Vector
Gaussian
Multiple
Description
Coding

$$\Theta_{1,1}^* \mathbf{M}_{1,1} = \mathbf{0}, \quad (1.4)$$

$$(\Theta_{2,1}^* - \Theta_{1,1}^*) \mathbf{M}_{2,1} = \mathbf{0}, \quad (\Theta_{2,2}^* - \Theta_{1,1}^*) \mathbf{M}_{2,2} = \mathbf{0}, \quad (1.5)$$

$$(\Sigma_X - \Theta_{2,1}^*) \mathbf{M}_{3,1} = \mathbf{0}, \quad (\Sigma_X - \Theta_{2,2}^*) \mathbf{M}_{3,2} = \mathbf{0}. \quad (1.6)$$

Problem
Formulation
A Single-Letter
Lower Bound for
Sum Rate
Vector Gaussian
Case

Vector
Gaussian
One-Help-One
Problem

Problem
Formulation
KKT Conditions
The Extremal
Inequality
Rate Region

Vector Gaussian Case: Direct Theorem

Two Results
on Vector
Gaussian
Multi-
Terminal
Source Coding

Yinfei Xu

Vector
Gaussian
Multiple
Description
Coding

Problem
Formulation
A Single-Letter
Lower Bound for
Sum Rate
Vector Gaussian
Case

Vector
Gaussian
One-Help-One
Problem

Problem
Formulation
KKT Conditions
The Extremal
Inequality
Rate Region

Distortion enhancement: define

$$(\Theta_{1,1}^* + \tilde{\Sigma}_{S_{1,1}})^{-1} = (\Theta_{1,1}^* + \Sigma_{S_{1,1}})^{-1} + \mathbf{M}_{1,1}$$

$$(\Theta_{1,1}^* + \tilde{\Sigma}_{S_{2,1}})^{-1} = (\Theta_{1,1}^* + \Sigma_{S_{2,1}})^{-1} + \mathbf{M}_{2,1}$$

$$(\Theta_{1,1}^* + \tilde{\Sigma}_{S_{2,2}})^{-1} = (\Theta_{1,1}^* + \Sigma_{S_{2,2}})^{-1} + \mathbf{M}_{2,2}$$

$$(\Theta_{2,1}^* + \tilde{\Sigma}_{S_{3,1}})^{-1} = (\Theta_{2,1}^* + \Sigma_{S_{3,1}})^{-1} + \mathbf{M}_{3,1}$$

$$(\Theta_{2,2}^* + \tilde{\Sigma}_{S_{3,3}})^{-1} = (\Theta_{2,2}^* + \Sigma_{S_{3,3}})^{-1} + \mathbf{M}_{3,2}$$

Along the same line of **channel enhancement** argument (H. Weingarten et al), the optimization lower bound can be written in another way.

Vector Gaussian Case: Direct Theorem

Two Results
on Vector
Gaussian
Multi-
Terminal
Source Coding

Yinfei Xu

Vector
Gaussian
Multiple
Description
Coding

Problem
Formulation
A Single-Letter
Lower Bound for
Sum Rate
Vector Gaussian
Case

Vector
Gaussian
One-Help-One
Problem

Problem
Formulation
KKT Conditions
The Extremal
Inequality
Rate Region

$$\begin{aligned} R(\underline{d}) &\geq \frac{1}{2} \log \frac{|\Sigma_X + \Sigma_{\tilde{S}_{1,1}}|}{|\Sigma_{\tilde{S}_{1,1}}|} \\ &+ \frac{1}{2} \log \frac{|\Theta_{1,1}^* + \Sigma_{\tilde{S}_{1,1}}| |\Sigma_X + \Sigma_{\tilde{S}_{2,1}}| |\Sigma_X + \Sigma_{\tilde{S}_{2,2}}|}{|\Sigma_X + \Sigma_{\tilde{S}_{1,1}}| |\Theta_{1,1}^* + \Sigma_{\tilde{S}_{2,1}}| |\Theta_{1,1}^* + \Sigma_{\tilde{S}_{2,2}}|} \\ &+ \frac{1}{2} \log \frac{|\Theta_{2,1}^* + \Sigma_{\tilde{S}_{2,1}}| |\Sigma_X + \Sigma_{\tilde{S}_{3,1}}| |\Sigma_X + \Sigma_{S_{3,2}}|}{|\Sigma_X + \Sigma_{\tilde{S}_{2,1}}| |\Theta_{2,1}^* + \Sigma_{\tilde{S}_{3,1}}| |\Theta_{2,1}^* + \Sigma_{S_{3,2}}|} \\ &+ \frac{1}{2} \log \frac{|\Theta_{2,2}^* + \Sigma_{\tilde{S}_{2,2}}| |\Sigma_X + \Sigma_{\tilde{S}_{3,3}}| |\Sigma_X + \Sigma_{S_{3,4}}|}{|\Sigma_X + \Sigma_{\tilde{S}_{2,2}}| |\Theta_{2,2}^* + \Sigma_{\tilde{S}_{3,3}}| |\Theta_{2,2}^* + \Sigma_{S_{3,4}}|} \end{aligned}$$

Vector Gaussian Case: Direct Theorem

Two Results
on Vector
Gaussian
Multi-
Terminal
Source Coding

Yinfei Xu

Vector
Gaussian
Multiple
Description
Coding

Problem
Formulation
A Single-Letter
Lower Bound for
Sum Rate
Vector Gaussian
Case

Vector
Gaussian
One-Help-One
Problem

Problem
Formulation
KKT Conditions
The Extremal
Inequality
Rate Region

$$(\Theta_{1,1}^* + \tilde{\Sigma}_{S_{2,1}})^{-1} + (\Theta_{1,1}^* + \tilde{\Sigma}_{S_{2,2}})^{-1} = (\Theta_{1,1}^* + \tilde{\Sigma}_{S_{1,1}})^{-1} \quad (1.7)$$

$$(\Theta_{2,1}^* + \tilde{\Sigma}_{S_{3,1}})^{-1} + (\Theta_{2,1}^* + \Sigma_{S_{3,2}})^{-1} = (\Theta_{2,1}^* + \tilde{\Sigma}_{S_{2,1}})^{-1} \quad (1.8)$$

$$(\Theta_{2,2}^* + \tilde{\Sigma}_{S_{3,3}})^{-1} + (\Theta_{2,2}^* + \Sigma_{S_{3,4}})^{-1} = (\Theta_{2,2}^* + \tilde{\Sigma}_{S_{2,2}})^{-1}, \quad (1.9)$$

Vector Gaussian Case: Direct Theorem

Two Results
on Vector
Gaussian
Multi-
Terminal
Source Coding

Yinfei Xu

Vector
Gaussian
Multiple
Description
Coding

Problem
Formulation
A Single-Letter
Lower Bound for
Sum Rate
Vector Gaussian
Case

Vector
Gaussian
One-Help-One
Problem

Problem
Formulation
KKT Conditions
The Extremal
Inequality
Rate Region

Define **independent** zero means Gaussian random vectors $W_{1,1}, (W_{2,1}, W_{2,2}), (W_{3,1}, W_{3,2}), (W_{3,3}, W_{3,4})$ with covariance:

$$\begin{aligned}\text{cov}(W_{1,1}) &= \tilde{\Sigma}_{S_{1,1}} \\ \text{cov}(W_{2,1}, W_{2,2}) &= \begin{pmatrix} \tilde{\Sigma}_{S_{2,1}} - \tilde{\Sigma}_{S_{1,1}} & -\Theta_{1,1}^* - \tilde{\Sigma}_{S_{1,1}} \\ -\Theta_{1,1}^* - \tilde{\Sigma}_{S_{1,1}} & \tilde{\Sigma}_{S_{2,2}} - \tilde{\Sigma}_{S_{1,1}} \end{pmatrix} \\ \text{cov}(W_{3,2i-1}, W_{3,2i}) &= \begin{pmatrix} \tilde{\Sigma}_{S_{3,2i-1}} - \tilde{\Sigma}_{S_{2,i}} & -\Theta_{2,i}^* - \tilde{\Sigma}_{S_{2,i}} \\ -\Theta_{2,i}^* - \tilde{\Sigma}_{S_{2,i}} & \tilde{\Sigma}_{S_{3,2i}} - \tilde{\Sigma}_{S_{3,i}} \end{pmatrix} \\ & \quad i = 1, 2.\end{aligned}$$

Vector Gaussian Case: Direct Theorem

Theorem

For $k = 1, 2, 3; i = 1, \dots, 2^{k-1}$,

- $\text{cov}(W_{k,2i-1}, W_{k,2i}) \succeq \mathbf{0}$,
- *There exists $(\mathbf{H}_{k+1,2i-1}, \mathbf{H}_{k+1,2i})$ such that the following identities hold*

$$\mathbf{H}_{k+1,2i-1} + \mathbf{H}_{k+1,2i} = \mathbf{I}$$

$$\mathbf{H}_{k+1,2i-1} W_{k+1,2i-1} + \mathbf{H}_{k+1,2i} W_{k+1,2i} = \mathbf{0}.$$

Define $U_j, j = 1, 2, 3, 4$ as output of each description

$$U_1 = X + W_{1,1} + W_{2,1} + W_{3,1},$$

$$U_2 = X + W_{1,1} + W_{2,1} + W_{3,2},$$

$$U_3 = X + W_{1,1} + W_{2,2} + W_{3,3},$$

$$U_4 = X + W_{1,1} + W_{2,2} + W_{3,4}.$$

Vector Gaussian Case: Direct Theorem

Two Results
on Vector
Gaussian
Multi-
Terminal
Source Coding

Yinfei Xu

Vector
Gaussian
Multiple
Description
Coding
Problem
Formulation
A Single-Letter
Lower Bound for
Sum Rate
Vector Gaussian
Case

Vector
Gaussian
One-Help-One
Problem

Problem
Formulation
KKT Conditions
The Extremal
Inequality
Rate Region

Check the covariance: $\text{cov}(X|U_{k,i})$, $k = 1, 2, 3; i = 1, \dots, 2^{k-1}$.

e.g.

$$\begin{aligned} & \text{cov}(X|U_1, U_2) \\ &= \text{cov}(X|X + W_{1,1} + W_{2,1} + W_{3,1}, X + W_{1,1} + W_{2,1} + W_{3,2}) \\ &= \text{cov}(X|\mathbf{H}_{3,1}(X + W_{1,1} + W_{2,1} + W_{3,1}) \\ & \quad + \mathbf{H}_{3,2}(X + W_{1,1} + W_{2,1} + W_{3,2})) \\ &= \text{cov}(X|X + W_{1,1} + W_{2,1}) \\ &= \left(\Sigma_X^{-1} + \tilde{\Sigma}_{S_{2,1}}^{-1} \right)^{-1} \preceq \left(\Sigma_X^{-1} + \Sigma_{S_{2,1}}^{-1} \right)^{-1} = \mathbf{D}_{S_{2,1}} \end{aligned}$$

Vector Gaussian Case: Direct Theorem

Two Results
on Vector
Gaussian
Multi-
Terminal
Source Coding

Yinfei Xu

Vector
Gaussian
Multiple
Description
Coding

Problem
Formulation
A Single-Letter
Lower Bound for
Sum Rate
Vector Gaussian
Case

Vector
Gaussian
One-Help-One
Problem

Problem
Formulation
KKT Conditions
The Extremal
Inequality
Rate Region

Check the sum rate for El Gamal-Cover inner bound

$$\begin{aligned} & h(U_1) + h(U_2) + h(U_3) + h(U_4) - h(U_1, U_2, U_3, U_4|X) \\ &= I(X; U_1, U_2, U_3, U_4) + I(U_1, U_2; U_3, U_4) + \\ & \quad I(U_1; U_2) + I(U_3; U_4) \end{aligned}$$

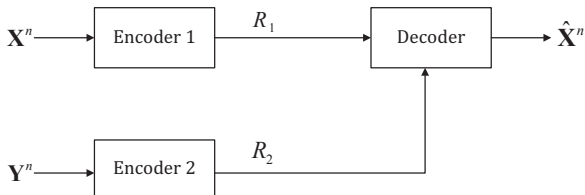
Vector Gaussian Case: Direct Theorem

Check the mutual information term: e.g. $I(U_1; U_2)$:

Write $X = \tilde{X} + \hat{X}$, where \tilde{X} and \hat{X} have covariance $\Sigma_X - \Theta_{2,1}^*$ and $\Theta_{2,1}^*$.

$$\begin{aligned} & I(U_1; U_2) \\ &= I(\tilde{X}; U_1) + I(\tilde{X}; U_2) - I(\tilde{X}; U_1, U_2) + I(U_1; U_2 | \tilde{X}) \\ &= I(\tilde{X}; \tilde{X} + \hat{X} + W_{1,1} + W_{2,1} + W_{3,1}) \left(\frac{1}{2} \log \frac{|\Sigma_X + \Sigma_{\tilde{S}_{3,1}}|}{|\Theta_{2,1}^* + \Sigma_{\tilde{S}_{3,1}}|} \right) \\ &+ I(\tilde{X}; \tilde{X} + \hat{X} + W_{1,1} + W_{2,1} + W_{3,2}) \left(\frac{1}{2} \log \frac{|\Sigma_X + \Sigma_{\tilde{S}_{3,2}}|}{|\Theta_{2,1}^* + \Sigma_{\tilde{S}_{3,2}}|} \right) \\ &- I(\tilde{X}; \tilde{X} + \hat{X} + W_{1,1} + W_{2,1}) \left(-\frac{1}{2} \log \frac{|\Sigma_X + \Sigma_{\tilde{S}_{2,1}}|}{|\Theta_{2,1}^* + \Sigma_{\tilde{S}_{2,1}}|} \right) \end{aligned}$$

Problem Formulation

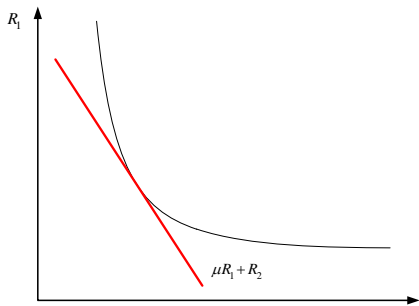


- $\{\mathbf{X}_i, \mathbf{Y}_i\}_{i=1}^{\infty}$ be a sequence of i.i.d. random vectors.
- $\mathbf{X}_i \sim \mathcal{N}(\mathbf{0}, \mathbf{K}_X)$; $\mathbf{Y}_i \sim \mathcal{N}(\mathbf{0}, \mathbf{K}_Y)$.
- **Correlation:** $\mathbf{X} = \mathbf{Y} + \mathbf{N}$, $\mathbf{N} \sim \mathcal{N}(\mathbf{0}, \mathbf{K}_N)$.

Problem Formulation

- The **rate-distortion region** $\mathcal{R}(\mathbf{D})$ is the closure of all achievable rate tuples (R_1, R_2) subject to distortion \mathbf{D} .
- Tangent hyperplane characterization: given any fixed $\mathbf{D} \succ \mathbf{0}$ and $\mu \geq 0$,

$$R(\mathbf{D}, \mu) = \inf_{(R_1, R_2) \in \mathcal{R}(\mathbf{D})} \mu R_1 + R_2$$



Problem Formulation

Two Results
on Vector
Gaussian
Multi-
Terminal
Source Coding

Yinfei Xu

Vector
Gaussian
Multiple
Description
Coding

Problem
Formulation
A Single-Letter
Lower Bound for
Sum Rate
Vector Gaussian
Case

Vector
Gaussian
One-Help-One
Problem

Problem
Formulation
KKT Conditions
The Extremal
Inequality
Rate Region

- Berger-Tung inner bound:

$$R^*(\mathbf{D}, \mu) \triangleq \min_{\mathbf{B}_1, \mathbf{B}_2} \frac{\mu}{2} \log \frac{|\mathbf{K}_X - \mathbf{B}_2|}{|\mathbf{K}_X - \mathbf{B}_1 - \mathbf{B}_2|} + \frac{1}{2} \log \frac{|\mathbf{K}_Y|}{|\mathbf{K}_Y - \mathbf{B}_2|}$$

subject to $\mathbf{B}_1, \mathbf{B}_2 \succcurlyeq \mathbf{0}$, and
 $\mathbf{D} \succcurlyeq \mathbf{K}_X - \mathbf{B}_1 - \mathbf{B}_2$.

- Rahman and Wagner's theorem: $R(\mathbf{D}, \mu) = R^*(\mathbf{D}, \mu)$.
"distortion projection" + "source enhancement"
- Our proof: **Perturbation**
(An enhanced Liu-Viswanath's extremal inequality)

KKT Conditions for Berger-Tung inner Bound

Two Results
on Vector
Gaussian
Multi-
Terminal
Source Coding

Yinfei Xu

Vector
Gaussian
Multiple
Description
Coding

Problem
Formulation
A Single-Letter
Lower Bound for
Sum Rate
Vector Gaussian
Case

Vector
Gaussian
One-Help-One
Problem

Problem
Formulation
KKT Conditions
The Extremal
Inequality
Rate Region

$$\frac{\mu}{2}(\mathbf{K}_X - \mathbf{B}_1^* - \mathbf{B}_2^*)^{-1} - \Psi_1 - \Lambda = \mathbf{0}, \quad (2.1)$$

$$\begin{aligned} \frac{\mu}{2}(\mathbf{K}_X - \mathbf{B}_1^* - \mathbf{B}_2^*)^{-1} - \frac{\mu}{2}(\mathbf{K}_X - \mathbf{B}_2^*)^{-1} \\ + \frac{1}{2}(\mathbf{K}_Y - \mathbf{B}_2^*)^{-1} - \Psi_2 - \Lambda = \mathbf{0}, \end{aligned} \quad (2.2)$$

$$\mathbf{B}_1^* \Psi_1 = \mathbf{0}, \quad (2.3)$$

$$\mathbf{B}_2^* \Psi_2 = \mathbf{0}, \quad (2.4)$$

$$(\mathbf{D} - \mathbf{K}_X + \mathbf{B}_1^* + \mathbf{B}_2^*) \Lambda = \mathbf{0}, \quad (2.5)$$

$$\Psi_1, \Psi_2, \Lambda \succeq \mathbf{0}.$$

KKT Conditions for Berger-Tung inner Bound

Two Results
on Vector
Gaussian
Multi-
Terminal
Source Coding

Yinfei Xu

Vector
Gaussian
Multiple
Description
Coding

Problem
Formulation
A Single-Letter
Lower Bound for
Sum Rate
Vector Gaussian
Case

Vector
Gaussian
One-Help-One
Problem

Problem
Formulation
KKT Conditions
The Extremal
Inequality
Rate Region

$$\frac{\mu}{2}(\mathbf{K}_X - \mathbf{B}_1^* - \mathbf{B}_2^*)^{-1} = \Psi_1 + \Lambda, \quad (2.1)$$

$$\frac{\mu}{2}(\mathbf{K}_X - \mathbf{B}_2^*)^{-1} - \frac{1}{2}(\mathbf{K}_Y - \mathbf{B}_2^*)^{-1} = \Psi_2 + \Lambda, \quad (2.2)$$

$$\mathbf{B}_1^* \Psi_1 = \mathbf{0}, \quad (2.3)$$

$$\mathbf{B}_2^* \Psi_2 = \mathbf{0}, \quad (2.4)$$

$$(\mathbf{D} - \mathbf{K}_X + \mathbf{B}_1^* + \mathbf{B}_2^*)\Lambda = \mathbf{0}, \quad (2.5)$$

$$\Psi_1, \Psi_2, \Lambda \succeq \mathbf{0}.$$

The Extremal Inequality

Two Results
on Vector
Gaussian
Multi-
Terminal
Source Coding

Yinfei Xu

Vector
Gaussian
Multiple
Description
Coding

Problem
Formulation
A Single-Letter
Lower Bound for
Sum Rate
Vector Gaussian
Case

Vector
Gaussian
One-Help-One
Problem

Problem
Formulation
KKT Conditions
The Extremal
Inequality
Rate Region

Theorem

Let $(\mathbf{B}_1^*, \mathbf{B}_2^*)$ be the optimal solution for the optimization problem $R^*(\mathbf{D}, \mu)$. Let V be a random variable such that $\mathbf{X} \rightarrow \mathbf{Y} \rightarrow V$ forms a Markov chain. Then for every distribution of (U, V) satisfying constraints:

$$E \left[(\mathbf{X} - E[\mathbf{X}|U, V]) (\mathbf{X} - E[\mathbf{X}|U, V])^T \right] \preceq \mathbf{D},$$

we have

$$\begin{aligned} & \mu h(\mathbf{X}|V) - h(\mathbf{Y}|V) - \mu h(\mathbf{X}|U, V) \\ & \geq \frac{\mu}{2} \log |(2\pi e)(\mathbf{K}_{\mathbf{X}} - \mathbf{B}_2^*)| - \frac{1}{2} \log |(2\pi e)(\mathbf{K}_{\mathbf{Y}} - \mathbf{B}_2^*)| \\ & \quad - \frac{\mu}{2} \log |(2\pi e)(\mathbf{K}_{\mathbf{X}} - \mathbf{B}_2^* - \mathbf{B}_1^*)|. \end{aligned}$$

The Extremal Inequality

Two Results
on Vector
Gaussian
Multi-
Terminal
Source Coding

Yinfei Xu

Vector
Gaussian
Multiple
Description
Coding

Problem
Formulation
A Single-Letter
Lower Bound for
Sum Rate
Vector Gaussian
Case

Vector
Gaussian
One-Help-One
Problem

Problem
Formulation
KKT Conditions
The Extremal
Inequality
Rate Region

- **Covariance preserving transform:** for any $\gamma \in (0, 1)$,

$$\mathbf{X}_{1,\gamma} = \sqrt{1-\gamma}\mathbf{X} + \sqrt{\gamma}\mathbf{X}_1^G, \quad \mathbf{X}_1^G \sim \mathcal{N}(\mathbf{0}, \mathbf{K}_X - \mathbf{B}_2^* - \mathbf{B}_1^*).$$

$$\mathbf{Y}_{2,\gamma} = \sqrt{1-\gamma}\mathbf{Y} + \sqrt{\gamma}\mathbf{Y}_2^G, \quad \mathbf{Y}_2^G \sim \mathcal{N}(\mathbf{0}, \mathbf{K}_Y - \mathbf{B}_2^*).$$

$$\mathbf{X}_{2,\gamma} = \sqrt{1-\gamma}\mathbf{X} + \sqrt{\gamma}\mathbf{X}_2^G, \quad \mathbf{X}_2^G \sim \mathcal{N}(\mathbf{0}, \mathbf{K}_X - \mathbf{B}_2^*).$$

- Degraded assumptions:

$$\mathbf{X}_2^G = \mathbf{Y}_2^G + \mathbf{N}^G, \quad \mathbf{N}^G \sim \mathcal{N}(\mathbf{0}, \mathbf{K}_N),$$

$$\mathbf{X}_2^G = \mathbf{X}_1^G + \mathbf{W}^G, \quad \mathbf{W}^G \sim \mathcal{N}(\mathbf{0}, \mathbf{B}_1^*)$$

- Define

$$g(\gamma) = \mu h(\mathbf{X}_{2,\gamma}|V) - h(\mathbf{Y}_{2,\gamma}|V) - \mu h(\mathbf{X}_{1,\gamma}|U, V).$$

The Extremal Inequality

Two Results
on Vector
Gaussian
Multi-
Terminal
Source Coding

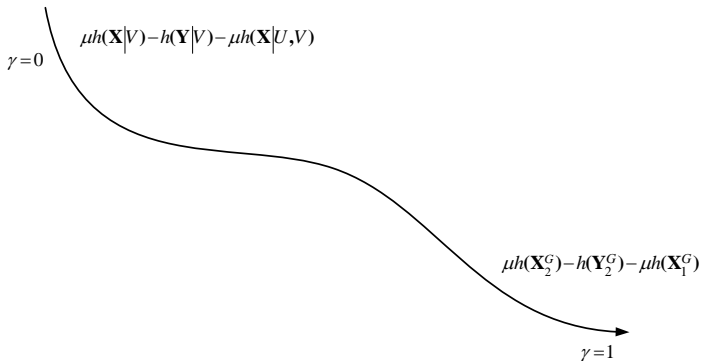
Yinfei Xu

Vector
Gaussian
Multiple
Description
Coding

Problem
Formulation
A Single-Letter
Lower Bound for
Sum Rate
Vector Gaussian
Case

Vector
Gaussian
One-Help-One
Problem

Problem
Formulation
KKT Conditions
The Extremal
Inequality
Rate Region



- If the inequality $(1 - \gamma)g'(\gamma) \leq 0$ is obtained, the extremal inequality is proved.

The Extremal Inequality

Two Results
on Vector
Gaussian
Multi-
Terminal
Source Coding

Yinfei Xu

Vector
Gaussian
Multiple
Description
Coding

Problem
Formulation
A Single-Letter
Lower Bound for
Sum Rate
Vector Gaussian
Case

Vector
Gaussian
One-Help-One
Problem

Problem
Formulation
KKT Conditions
The Extremal
Inequality
Rate Region

$$\begin{aligned} & 2(1 - \gamma)g'(\gamma) \\ = & \operatorname{tr} \left\{ \mu(\mathbf{K}_X - \mathbf{B}_2^*)^{-1} [(\mathbf{K}_X - \mathbf{B}_2^*)J(\mathbf{X}_{2,\gamma}|V)(\mathbf{K}_X - \mathbf{B}_2^*) \right. \\ & \quad \left. - (\mathbf{K}_X - \mathbf{B}_2^*)] \right. \\ & \quad \left. - (\mathbf{K}_Y - \mathbf{B}_2^*)^{-1} [(\mathbf{K}_Y - \mathbf{B}_2^*)J(\mathbf{Y}_{2,\gamma}|V)(\mathbf{K}_Y - \mathbf{B}_2^*) \right. \\ & \quad \left. - (\mathbf{K}_Y - \mathbf{B}_2^*)] \right. \\ & \quad \left. - \left\{ \mu(\mathbf{K}_X - \mathbf{B}_1^* - \mathbf{B}_2^*)^{-1} [(\mathbf{K}_X - \mathbf{B}_1^* - \mathbf{B}_2^*)J(\mathbf{X}_{1,\gamma}|U, V) \right. \right. \\ & \quad \left. \left. (\mathbf{K}_X - \mathbf{B}_1^* - \mathbf{B}_2^*) - (\mathbf{K}_X - \mathbf{B}_1^* - \mathbf{B}_2^*) \right\} \right\} \end{aligned}$$

Derivation of $g(\gamma)$ via **de Bruijn's Identity**

$$\frac{d}{d\gamma} h(\mathbf{X} + \sqrt{\gamma}\mathbf{N}|U) = \frac{1}{2} \operatorname{tr} \{ J(\mathbf{X} + \sqrt{\gamma}\mathbf{N}|U) \}$$

The Extremal Inequality

Two Results
on Vector
Gaussian
Multi-
Terminal
Source Coding

Yinfei Xu

Vector
Gaussian
Multiple
Description
Coding

Problem
Formulation
A Single-Letter
Lower Bound for
Sum Rate
Vector Gaussian
Case

Vector
Gaussian
One-Help-One
Problem

Problem
Formulation
KKT Conditions
The Extremal
Inequality
Rate Region

$$\begin{aligned} & 2(1 - \gamma)g'(\gamma) \\ \leq & \text{tr} \left\{ \mu(\mathbf{K}_X - \mathbf{B}_2^*)^{-1} [(\mathbf{K}_X - \mathbf{B}_2^*)J(\mathbf{X}_{2,\gamma}|V)(\mathbf{K}_X - \mathbf{B}_2^*) \right. \\ & \quad \left. - (\mathbf{K}_X - \mathbf{B}_2^*)] \right. \\ & \quad \left. - (\mathbf{K}_Y - \mathbf{B}_2^*)^{-1} [(\mathbf{K}_X - \mathbf{B}_2^*)J(\mathbf{X}_{2,\gamma}|V)(\mathbf{K}_X - \mathbf{B}_2^*) \right. \\ & \quad \left. - (\mathbf{K}_X - \mathbf{B}_2^*)] \right. \\ & \quad \left. - \left\{ \mu(\mathbf{K}_X - \mathbf{B}_1^* - \mathbf{B}_2^*)^{-1} [(\mathbf{K}_X - \mathbf{B}_1^* - \mathbf{B}_2^*)J(\mathbf{X}_{1,\gamma}|V) \right. \right. \\ & \quad \left. \left. (\mathbf{K}_X - \mathbf{B}_1^* - \mathbf{B}_2^*) - (\mathbf{K}_X - \mathbf{B}_1^* - \mathbf{B}_2^*) \right\} \right\} \end{aligned}$$

$$\mathbf{X}_{2,\gamma} \stackrel{(d)}{=} \tilde{\mathbf{X}}_{2,\gamma} = \mathbf{Y}_{2,\gamma} + \tilde{\mathbf{N}}, \text{ where } \tilde{\mathbf{N}} \stackrel{(d)}{=} \mathbf{N} \sim \mathcal{N}(\mathbf{0}, \mathbf{K}_N).$$

$$\begin{aligned} & (\mathbf{K}_Y - \mathbf{B}_2^*)J(\mathbf{Y}_{2,\gamma}|V)(\mathbf{K}_Y - \mathbf{B}_2^*) \\ \succeq & (\mathbf{K}_X - \mathbf{B}_2^*)J(\mathbf{X}_{2,\gamma}|V)(\mathbf{K}_X - \mathbf{B}_2^*) - \mathbf{K}_N. \end{aligned}$$

The Extremal Inequality

Two Results
on Vector
Gaussian
Multi-
Terminal
Source Coding

Yinfei Xu

Vector
Gaussian
Multiple
Description
Coding

Problem
Formulation
A Single-Letter
Lower Bound for
Sum Rate
Vector Gaussian
Case

Vector
Gaussian
One-Help-One
Problem

Problem
Formulation
KKT Conditions
The Extremal
Inequality
Rate Region

$$\begin{aligned} & (1 - \gamma)g'(\gamma) \\ \leq & \text{tr} \left\{ [(\mathbf{K}_X - \mathbf{B}_2^*)J(\mathbf{X}_{2,\gamma}|V)(\mathbf{K}_X - \mathbf{B}_2^*) - (\mathbf{K}_X - \mathbf{B}_2^* - \mathbf{B}_1^*) \right. \\ & \cdot J(\mathbf{X}_{1,\gamma}|U, V)(\mathbf{K}_X - \mathbf{B}_2^* - \mathbf{B}_1^*) + \mathbf{B}_1^*] \Psi_1 \left. \right\} \\ & - \text{tr} \left\{ [(\mathbf{K}_X - \mathbf{B}_2^*)J(\mathbf{X}_{2,\gamma}|V)(\mathbf{K}_X - \mathbf{B}_2^*) - (\mathbf{K}_X - \mathbf{B}_2^*)] \Psi_2 \right\} \\ & - \text{tr} \left\{ [(\mathbf{K}_X - \mathbf{B}_2^* - \mathbf{B}_1^*)J(\mathbf{X}_{1,\gamma}|U, V)(\mathbf{K}_X - \mathbf{B}_2^* - \mathbf{B}_1^*) \right. \\ & \left. - (\mathbf{K}_X - \mathbf{B}_2^* - \mathbf{B}_1^*)] \Lambda \right\} \end{aligned}$$

Substitute into KKT conditions (2.1) - (2.2),

$$\begin{aligned} \frac{\mu}{2}(\mathbf{K}_X - \mathbf{B}_1^* - \mathbf{B}_2^*)^{-1} &= \Psi_1 + \Lambda, \\ \frac{\mu}{2}(\mathbf{K}_X - \mathbf{B}_2^*)^{-1} - \frac{1}{2}(\mathbf{K}_Y - \mathbf{B}_2^*)^{-1} &= \Psi_2 + \Lambda, \end{aligned}$$

The Extremal Inequality

The first term:

- $\mathbf{X}_{2,\gamma} \stackrel{(d)}{=} \check{\mathbf{X}}_{2,\gamma} = \mathbf{X}_{1,\gamma} + \sqrt{\gamma} \check{\mathbf{W}}^G$, where $\check{\mathbf{W}}^G \sim \mathcal{N}(\mathbf{0}, \mathbf{B}_1^*)$.
- By **Fisher Information Matrix Inequality** and **Data Processing Inequality**:

$$\begin{aligned} & (\mathbf{K}_{\mathbf{X}} - \mathbf{B}_2^*) J(\mathbf{X}_{2,\gamma}|V) (\mathbf{K}_{\mathbf{X}} - \mathbf{B}_2^*) \\ & \preceq (\mathbf{K}_{\mathbf{X}} - \mathbf{B}_2^* - \mathbf{B}_1^*) J(\mathbf{X}_{1,\gamma}|V) (\mathbf{K}_{\mathbf{X}} - \mathbf{B}_2^* - \mathbf{B}_1^*) + \frac{1}{\gamma} \mathbf{B}_1^* \\ & \preceq (\mathbf{K}_{\mathbf{X}} - \mathbf{B}_2^* - \mathbf{B}_1^*) J(\mathbf{X}_{1,\gamma}|U, V) (\mathbf{K}_{\mathbf{X}} - \mathbf{B}_2^* - \mathbf{B}_1^*) + \frac{1}{\gamma} \mathbf{B}_1^*, \end{aligned}$$

- Therefore,

$$\begin{aligned} & \text{tr}\{[(\mathbf{K}_{\mathbf{X}} - \mathbf{B}_2^*) J(\mathbf{X}_{2,\gamma}|V) (\mathbf{K}_{\mathbf{X}} - \mathbf{B}_2^*) - (\mathbf{K}_{\mathbf{X}} - \mathbf{B}_2^* - \mathbf{B}_1^*) \\ & \quad \cdot J(\mathbf{X}_{1,\gamma}|U, V) (\mathbf{K}_{\mathbf{X}} - \mathbf{B}_2^* - \mathbf{B}_1^*) + \mathbf{B}_1^*] \Psi_1\} \\ & \leq \frac{1 + \gamma}{\gamma} \text{tr}\{\mathbf{B}_1^* \Psi_1\} = \mathbf{0}. \end{aligned}$$

The Extremal Inequality

Two Results
on Vector
Gaussian
Multi-
Terminal
Source Coding

Yinfei Xu

Vector
Gaussian
Multiple
Description
Coding

Problem
Formulation
A Single-Letter
Lower Bound for
Sum Rate
Vector Gaussian
Case

Vector
Gaussian
One-Help-One
Problem

Problem
Formulation
KKT Conditions
The Extremal
Inequality
Rate Region

The second term:

- By **Cramér–Rao Inequality**:

$$\begin{aligned} J(\mathbf{X}_{2,\gamma}|V)^{-1} &\preceq \text{cov}(\mathbf{X}_{2,\gamma}|V) \preceq \text{cov}(\mathbf{X}_{2,\gamma}) \\ &= \text{cov}(\sqrt{1-\gamma}\mathbf{X} + \sqrt{\gamma}\mathbf{X}_2^G) = \mathbf{K}_X - \gamma\mathbf{B}_2^*. \end{aligned}$$

- Therefore,

$$\begin{aligned} &\text{tr}\{[(\mathbf{K}_X - \mathbf{B}_2^*)J(\mathbf{X}_{2,\gamma}|V)(\mathbf{K}_X - \mathbf{B}_2^*) - (\mathbf{K}_X - \mathbf{B}_2^*)]\Psi_2\} \\ &\geq \text{tr}\{(\mathbf{K}_X - \mathbf{B}_2^*)(\mathbf{K}_X - \gamma\mathbf{B}_2^*)^{-1}(\gamma - 1)\mathbf{B}_2^*\Psi_2\} = 0. \end{aligned}$$

The Extremal Inequality

Two Results
on Vector
Gaussian
Multi-
Terminal
Source Coding

Yinfei Xu

Vector
Gaussian
Multiple
Description
Coding

Problem
Formulation
A Single-Letter
Lower Bound for
Sum Rate
Vector Gaussian
Case

Vector
Gaussian
One-Help-One
Problem

Problem
Formulation
KKT Conditions
The Extremal
Inequality
Rate Region

The third term:

- By **Cramér–Rao Inequality**:

$$J(\mathbf{X}_{1,\gamma}|U, V)^{-1} \succeq \text{cov}(\mathbf{X}_{1,\gamma}|U, V) \succeq \mathbf{D}.$$

- Therefore,

$$\begin{aligned} & \text{tr}\{[(\mathbf{K}_X - \mathbf{B}_2^* - \mathbf{B}_1^*)J(\mathbf{X}_{1,\gamma}|U, V)(\mathbf{K}_X - \mathbf{B}_2^* - \mathbf{B}_1^*) \\ & \quad - (\mathbf{K}_X - \mathbf{B}_2^* - \mathbf{B}_1^*)]\Lambda\} \\ & \geq \text{tr}\{(\mathbf{K}_X - \mathbf{B}_2^* - \mathbf{B}_1^*)\mathbf{D}^{-1}[\mathbf{K}_X - \mathbf{B}_2^* - \mathbf{B}_1^* - \mathbf{D}]\Lambda\} = 0. \end{aligned}$$

⇒ The proof of the extremal inequality is completed.

Rate Region

Two Results
on Vector
Gaussian
Multi-
Terminal
Source Coding

Yinfei Xu

Vector
Gaussian
Multiple
Description
Coding

Problem
Formulation
A Single-Letter
Lower Bound for
Sum Rate
Vector Gaussian
Case

Vector
Gaussian
One-Help-One
Problem

Problem
Formulation
KKT Conditions
The Extremal
Inequality
Rate Region

$$\begin{aligned}\mu R_1 + R_2 &\geq \mu I(\mathbf{X}; U|V) + I(\mathbf{Y}; V) \\ &= \mu h(\mathbf{X}|V) - \mu h(\mathbf{X}|U, V) - h(\mathbf{Y}|V) + h(\mathbf{Y}) \\ &\geq \frac{\mu}{2} \log \frac{|\mathbf{K}_X - \mathbf{B}_2^*|}{|\mathbf{K}_X - \mathbf{B}_1^* - \mathbf{B}_2^*|} + \frac{1}{2} \log \frac{|\mathbf{K}_Y|}{|\mathbf{K}_Y - \mathbf{B}_2^*|} \\ &= R^*(\mathbf{D}, \mu).\end{aligned}$$

Two Results
on Vector
Gaussian
Multi-
Terminal
Source Coding

Yinfei Xu

Vector
Gaussian
Multiple
Description
Coding
Problem
Formulation
A Single-Letter
Lower Bound for
Sum Rate
Vector Gaussian
Case

Vector
Gaussian
One-Help-One
Problem
Problem
Formulation
KKT Conditions
The Extremal
Inequality
Rate Region

Thank you!

Q & A ?